

The RUB logo consists of the letters 'RUB' in a bold, white, sans-serif font, centered within a dark blue rectangular background.

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On the Security of TLS Renegotiation

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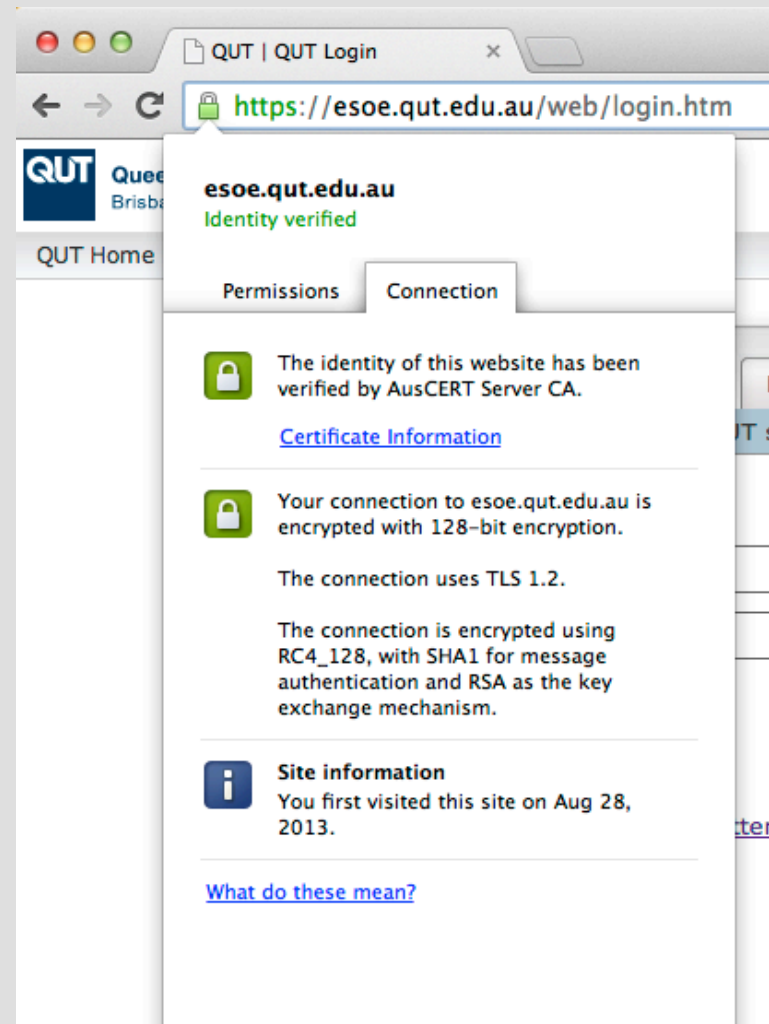
European Network of Excellence in
Cryptology II (ECRYPT II)

Agenda

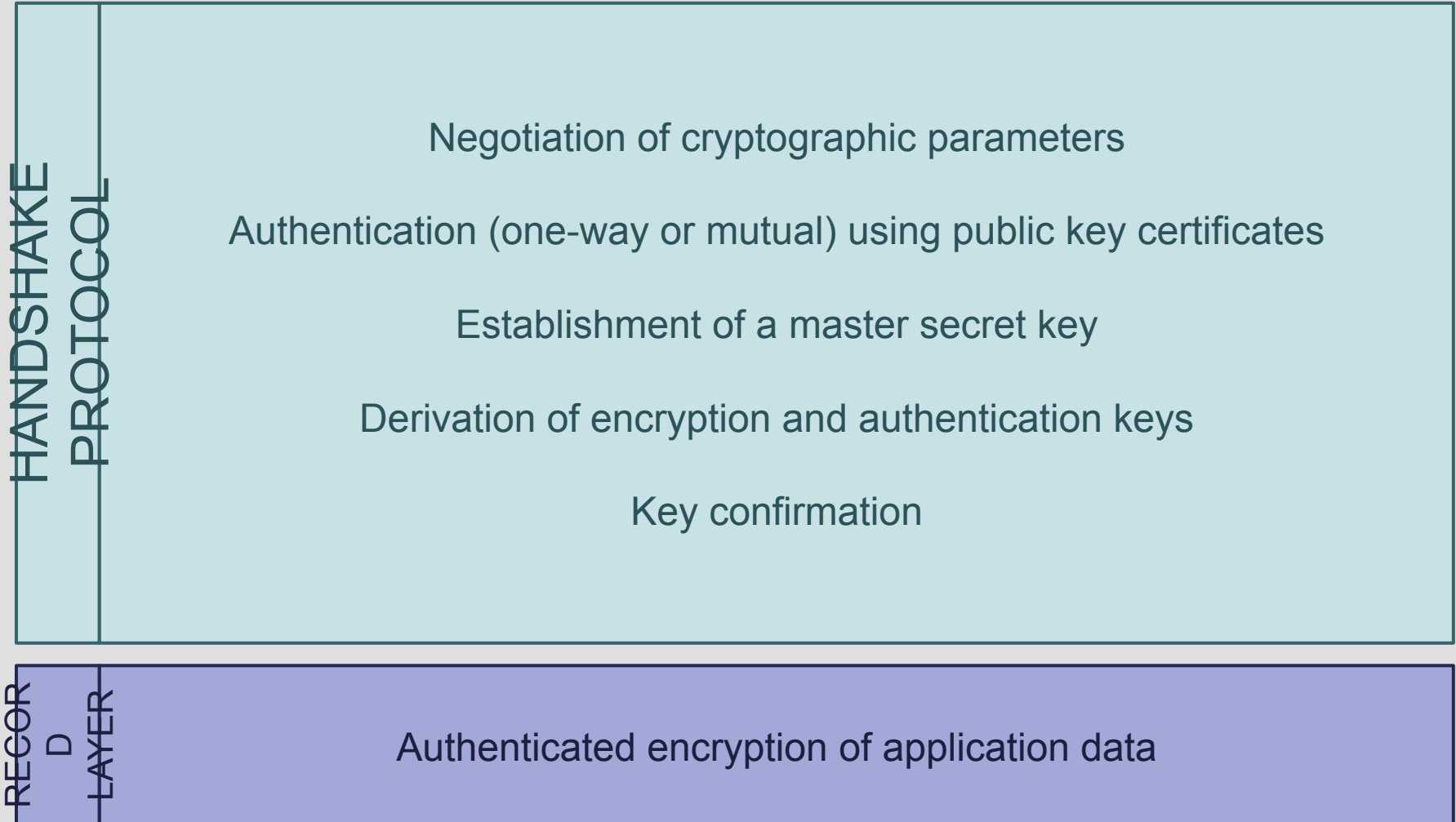
- What is TLS?
- TLS Renegotiation
- Renegotiation security

TLS (Transport Layer Security) protocol a.k.a. SSL (Secure Sockets Layer)

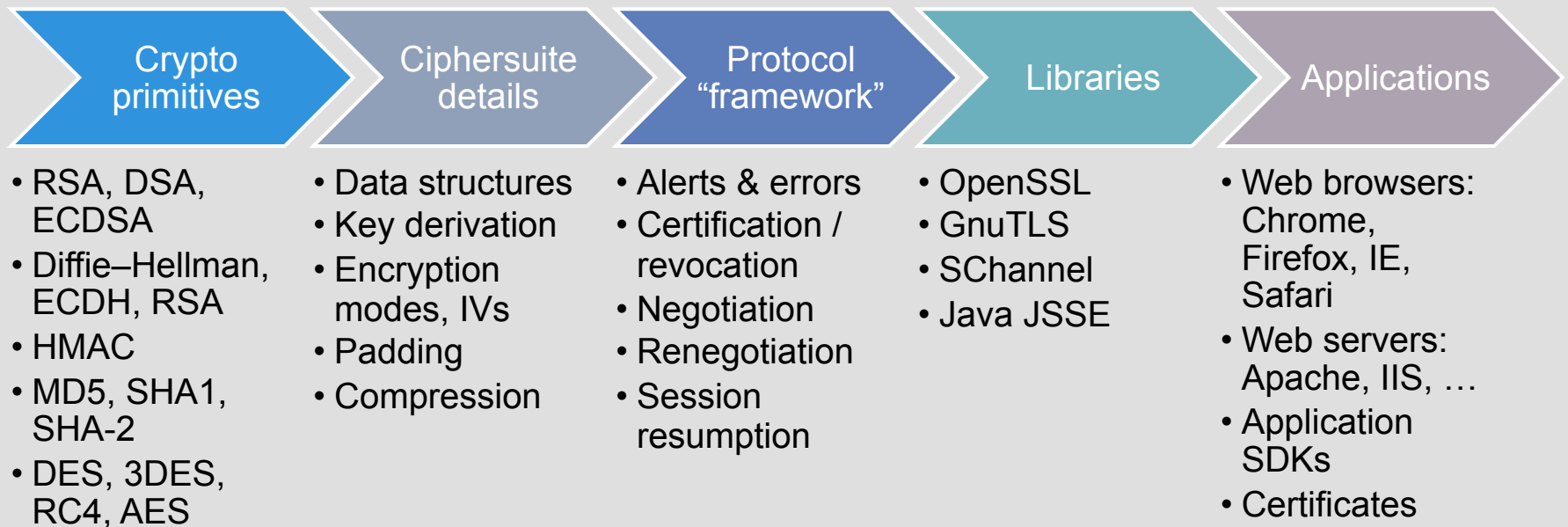
- Protects http, ftp, smtp, imap, ...
- The **most important cryptographic protocol on the Internet** — used to secure billions of connections every day.




Structure of TLS



TLS from primitive to application




Is TLS secure?



SSL v3.0
standardized

1996



Some variant
of one
ciphersuite of
the TLS
record layer
is a secure
encryption
scheme

2001

[Kra01]



Truncated
TLS
handshake
using RSA
key transport
is a secure
authenticated
key exchange
protocol

2002

[JK02]



Truncated
TLS
handshake
using RSA
key transport
or signed
Diffie–
Hellman is a
secure
authenticated
key exchange
protocol


2008

[MSW08]

“some variant”... “truncated TLS”...
limited ciphersuites

Is TLS secure?

1996  SSL v3.0 standardized

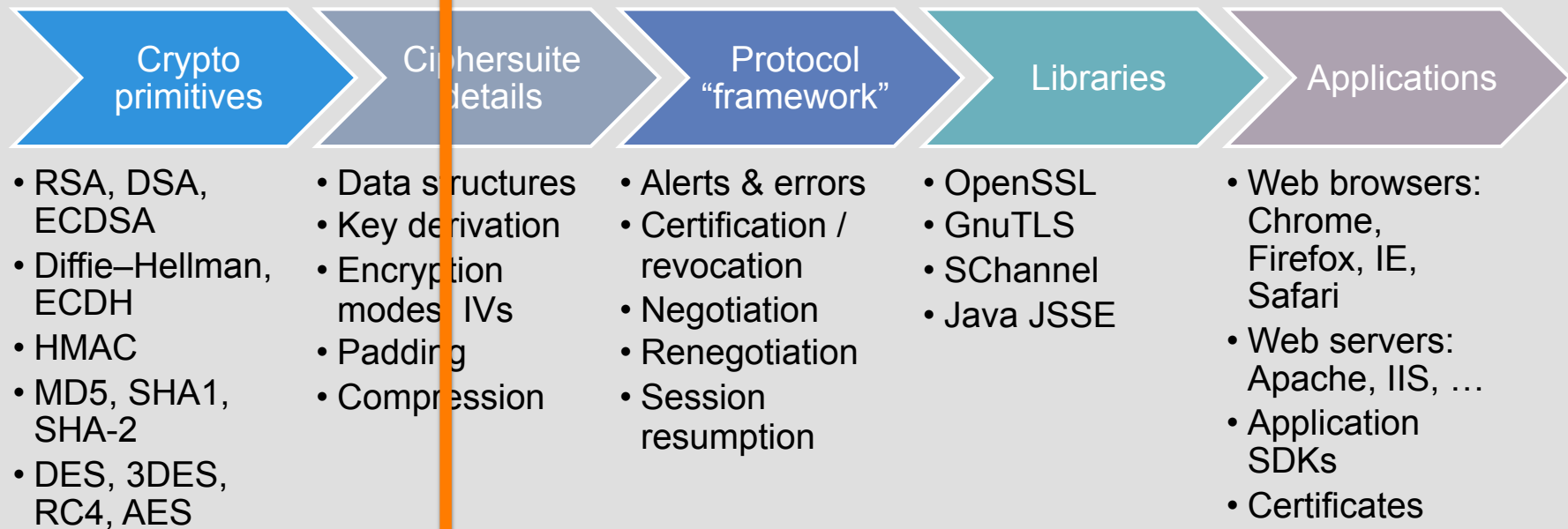
2011  Some modes of TLS record layer are secure authenticated encryption schemes
[PRS11]

2012  Unaltered full signed Diffie–Hellman ciphersuite provides a secure channel
[JKSS12]

2013  Most unaltered full TLS ciphersuites provide a secure channel
[KSS13, KPW13]

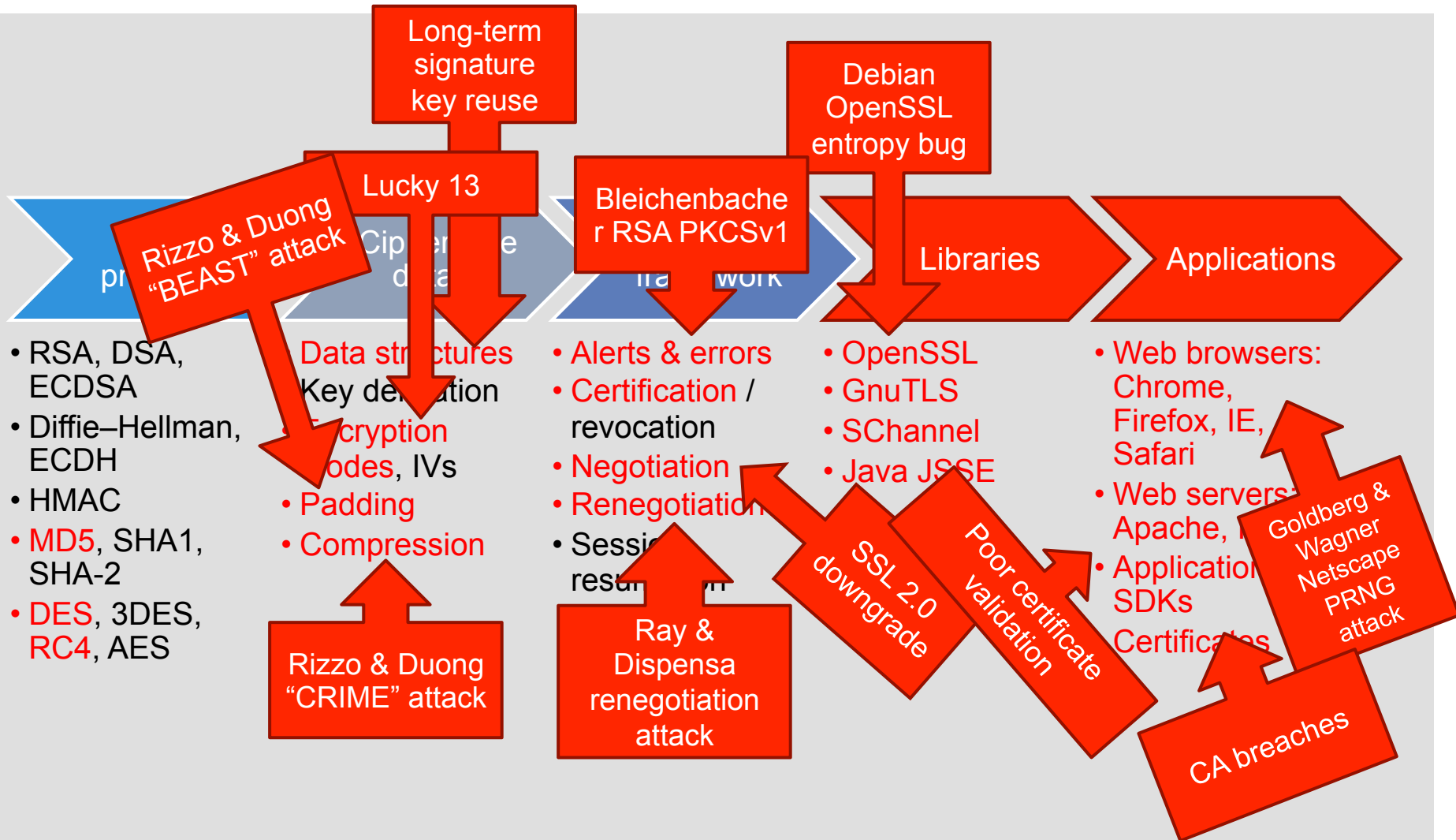
“unaltered”... “full”... “most ciphersuites”

Results on the provable security of TLS



Good combinations of these are theoretically secure, when done properly.

Real-world attacks on TLS



TLS is secure, but NOT

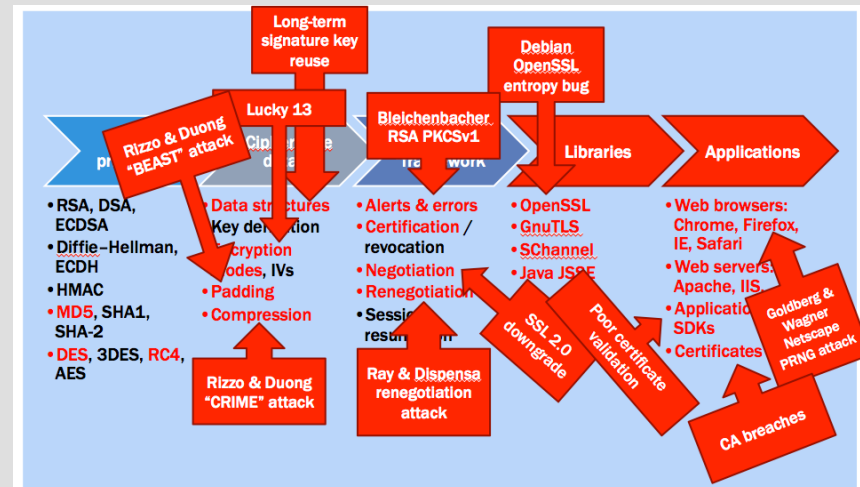
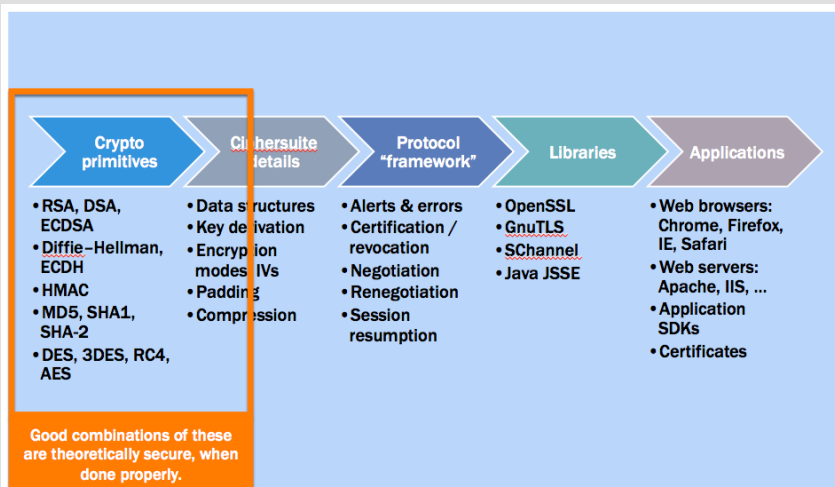
- TLS_DHE is a secure ACCE Protocol [JKSS12]
- Most TLS ciphersuits are secure ACCE Protocols [KSS13, KPW13]

- BEAST attack [DR11]
- Renegotiation attack on TLS [RD09]

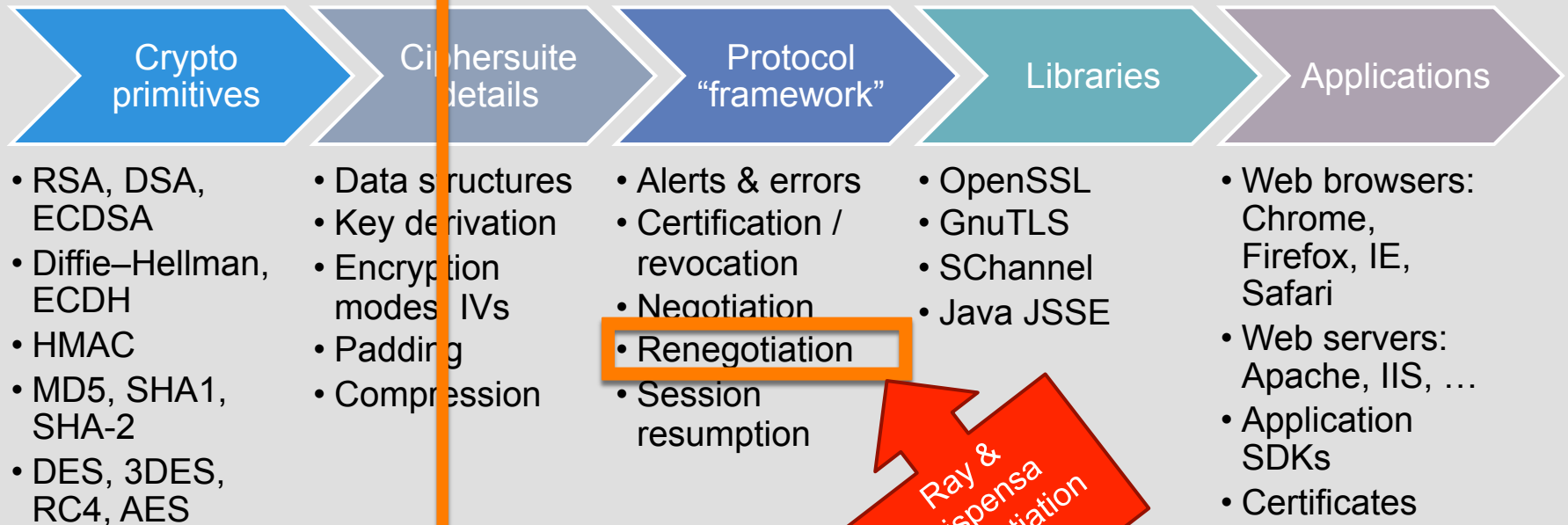
The gap between theory and practice

Provable security results

Real-world attacks



Our focus on TLS



Good combinations of these are theoretically secure, when done properly.

Renegotiation

Allows to create a new TLS channel that continues from the existing one.

Once you've established a TLS channel, why would you ever want to renegotiate it?

- Change cryptographic parameters
- Refresh encryption keys
- Change authentication credentials
- Identity hiding for client
 1. Establish a one-way authenticated TLS session
 2. Renegotiate using mutual authentication.
Since handshake messages are sent in the encrypted TLS channel, client's identity is kept private.

Renegotiation in TLS

(pre-November 2009)

Client

Server
(TLS)

TLS handshake₀

TLS recordlayer₀

m₀

I'd like to renegotiate

Messages are like those in original handshake, just encrypted

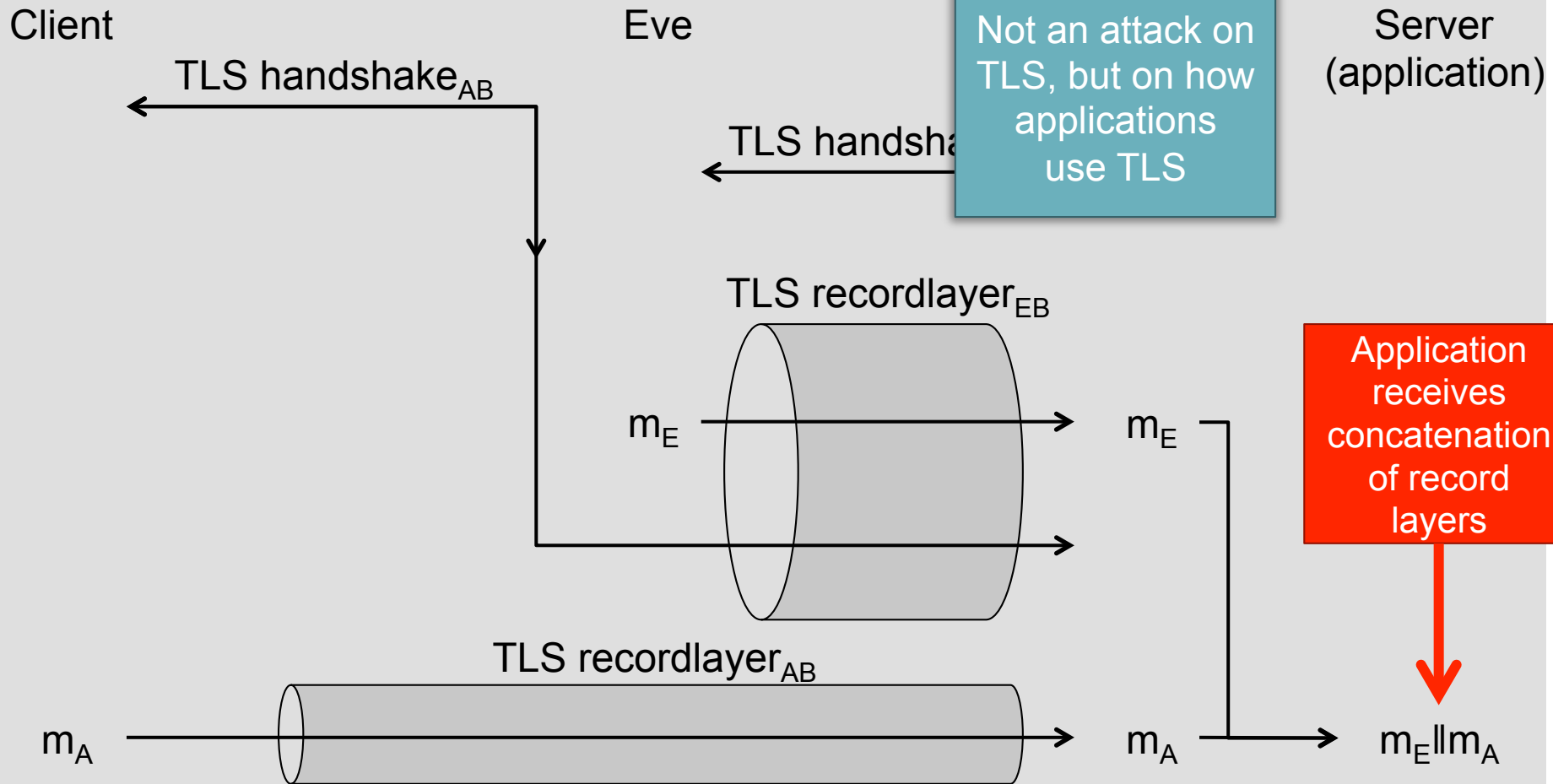
TLS handshake₁

TLS recordlayer₁

m₁

TLS Renegotiation “Attack”

Ray & Dispensa, November 2009



Eve orders a pizza but Alice pays

- Eve sends
 - $m_E = \text{"GET /orderPizza?deliverTo=123-Fake-St} \leftarrow \text{X-Ignore-This: "}$
- Client sends
 - $m_A = \text{"GET /orderPizza?deliverTo=456-Real-St} \leftarrow \text{Cookie: } \leftarrow \text{Account=1A2B"}$
- Server's web server receives
 - $m_E \parallel m_A = \text{"GET /orderPizza?deliverTo=123-Fake-St} \leftarrow \text{X-Ignore-This: GET /orderPizza?deliverTo=456-Real-St} \leftarrow \text{Cookie: Account=1A2B"}$
 - X-Ignore-This: is an invalid header, so the rest of that line gets ignored.
 - Eve's GET request is processed with the cookie supplied by the client.

Renegotiation security

- Q: What property should a secure renegotiable protocol intuitively have?
- A: Whenever two parties successfully renegotiate, they are assured they have the exact same view of everything that happened previously.

Technical approach

1. Extend authenticated and confidential channel establishment (ACCE) security model to include renegotiable, multi-phase protocols.
2. Define security notion for renegotiable protocols.
 - secure renegotiable ACCE (full)
 - weakly secure renegotiable ACCE
3. Show that TLS without fixes does not satisfy security definition.
4. Show that TLS with fixes does satisfy (weak) security definition.
 - Generic reduction: If a TLS ciphersuite satisfies a certain property, then, when combined with fixes, it is a weakly secure renegotiable ACCE.
5. Propose fix to achieve (full) security.

ACCE security

Authenticated and Confidential Channel Establishment

Adversary's goals:

1. Violate authentication:

- make Alice accept where her intended partner Bob is uncorrupted but has no matching conversation

2. Violate ciphertext integrity or confidentiality:

- distinguish which of two (chosen) plaintexts were encrypted

Provable security of TLS:

- TLS_DHE_DSS_WITH_3DES_EDE_CBC_SHA is a secure ACCE protocol
 - (JKSS; CRYPTO 2012)
- Most TLS ciphersuites are ACCE secure
 - (KSS; ePrint 2013, KPR; CRYPTO 2013)

Secure renegotiable ACCE

■ Secure renegotiable ACCE (full):

- Authentication:
 - when a party successfully renegotiate a new phase, its partner has a phase with a matching handshake and record layer transcript.

■ Weakly secure renegotiable ACCE:

- Authentication:
 - when a party successfully renegotiate a new phase, its partner has a phase with a matching handshake and record layer transcript, ***provided no previous phase's session key was revealed***

Generic renegotiation security of TLS

- Prove generically that any ACCE-secure TLS ciphersuit is also renegotiation-secure
- Can't do this because of Ray-Dispensa attack
- Next best thing: prove that any ACCE-secure TLS ciphersuite with the fixes is also renegotiation-secure
- Problem: hard to generically insert the fixes into existing ACCE protocols

Main technical theorem

- Suppose a TLS ciphersuite is ACCE-secure, even if
 - The Finished message is revealed
 - Arbitrary „tags“ are included in Client/ServerHello messages
- Then that ciphersuite is weakly renegotiable secure if the RFC5746 fixes are used.

Weakly secure renegotiable ACCE

Definition

- Weakly secure renegotiable ACCE:
 - Authentication:
 - when a party successfully renegotiate a new phase, its partner has a phase with a matching handshake and record layer transcript, *provided no previous phase's session key was revealed*

TLS

- TLS without fixes is not a weakly secure renegotiable ACCE.
- TLS with RFC 5746 fixes is a weakly secure renegotiable ACCE.

Secure renegotiable ACCE

Definition

- Secure renegotiable ACCE (full):
 - Authentication:
 - when a party successfully renegotiate a new phase, its partner has a phase with a matching handshake and record layer transcript, *provided no previous phase's session key was revealed.*

TLS

- TLS with or without RFC 5746 fixes is not a secure renegotiable ACCE.
- Can be fixed by including hash of previous record layer messages in RIE.

Conclusion

Theory

- First paper investigated renegotiation for cryptographic protocols.
 - Different levels of renegotiation security
- Security of a protocol in isolation doesn't imply security with renegotiation.

Practice

- The standardized fixes are provably secure.
- Safe to use fixed renegotiation in TLS.

Thank you very much for your attention

Now

Later

Questions?

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